Extraction of drainage network from grid terrain datasets using parallel computing

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Abstract—Drainage-network extraction from grid terrain datasets is the essential content for hydrologic analysis, soil erosion and geomorphology. However, drainage-network extraction is normally very time-consuming using the sequential program, especially for grid terrain datasets with high resolution and large scope. This paper proposed a set of parallel algorithms to preprocess DEM, determine flow routing and delineate drainage network using cluster computers in an MPI programming model. Both depressions and flat areas could be processed by the DEM preprocessing algorithm. For using the number of upslopedependence neighbors of a cell, the flow-routing algorithm makes flow accumulation calculation similar to a local algorithm. The experiment results shows that the proposed parallel approach to extract drainage-network performs much better than sequential algorithms and has the better parallel efficiency.

I. INTRODUCTION

As one of the most significant spatial data, digital elevation models (DEMs) are data structures representing rectangular grids of terrain data composed of cells arranged as a raster, where the value of each cell denotes the elevation of that geographic point above some base value [1]. Based on DEMs, it could mine much useful information about the morphology of a terrain surface via the technology of digital terrain analysis (DTA). The drainage network derived from DEMs is one of the key contents in DTA and has a marked influence on many fields, such as hydrologic analysis, soil erosion and geomorphology [2, 3]. Commonly, drainage-network delineation from DEMs is based on a model for representation of flow routing across the surface and between grid cells [4]. To achieve flow routing, it need to firstly determine a flow direction to each grid cell, and then calculate flow accumulations. According to the flow allocation ratio from a cell when it drains into the neighboring cells, existing flow-direction algorithms can be classified into two main types: single-flowdirection (SFD) algorithms [4, 5] and multiple-flow-direction (MFD) algorithms [6, 7]. Among all the flow-direction algorithms, the D8 algorithm proposed by O'Callaghan and Mark (1984) is widely applied for its simplicity and efficiency [8,9].

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From the flow direction grid, a flow accumulation grid can be achieved, and then drainage network can be extracted by means of a threshold [4].

During drainage-network extraction for real DEMs, a DEM preprocessing algorithm is normally applied to fill in the depressions and remove the flat areas in the DEMs before calculating the flow direction [2]. These depressions and flat areas which are both from real topographic features and errors introduced during the DEM production process would cause to fail to obtain correct result for flow direction. Over the last twenty years, many DEM preprocessing algorithms have been develop to obtain a smooth DEM without any depression and flat area [10, 11]. The algorithm proposed in [11] has generally been recognized to perform better from the perspective of algorithm efficiency [12].

At present, along with the development of spatial data capture technology, DEMs have increasingly been big datasets for large scopes and high resolutions [13]. Meanwhile, drainage-network extraction is normally very time-consuming caused not only from the recursive routing algorithm, but also from the iterative DEM preprocessing algorithm. As a result, the traditional serial algorithm for calculating drainage basin is usually hard to meet users' demand of time responding for its long execution time. To speed up the execution time, some researchers have proposed parallel solution for drainage-network extraction [14-16]. These researches focused on the segment of drainage-network extraction, however, DEM preprocessing phase, by far the correction procedure and time consuming, was not processed in parallel. Wallis [17] and Song [18] proposed a parallel DEM preprocessing algorithm from the method of Planchon and Darboux, respectively. However, these parallel algorithms could only fill the sinks to flat areas, and could not remove the flat areas with a very small slope gradient.

This paper presents a design and implementation of parallelized drainage-network extraction. In this work, it adopts Planchon and Darboux's method for DEM preprocessing and the

D8 approach for flow routing to parallelize algorithms using MPI. The remainder of this paper is organized as follows. In section 2, we introduce the implementations of both serial algorithms and parallel ones. In section 3, the performance results from parallel computing are compared with sequential computing on a small cluster, and in Section 4, the main conclusions are drawn.

II. COMPUTATION ALGORITHMS FOR EXTRACTING DRAINAGE NETWORK

The drainage network is a basic feature of terrain surface. Extracting drainage network from DEMs commonly contains three steps: (1) DEM preprocessing, which fills dispersions and removes flat areas, (2) flow routing, which determines the flow direction and flow accumulation for each grid cell, and (3) drainage-network delineating, which identifies a grid cell as being part of a channel using a fixed threshold. Among them, the steps of DEM preprocessing and flow routing are main parts for drainage-network extraction. They are usually the iteration processes and very time consuming. When the first two steps implemented, a common approach for the last step is to decide a cell if its flow accumulation is greater than a fixed threshold [4].

A. DEM preprocessing

The DEM preprocessing algorithm proposed by Planchon and Darboux (called P&D algorithm for short) includes two stages: (1) the water-covering stage, which inundates a thick layer of water over the entire DEM except for the boundary, namely cells on boundary initialized values of the original DEM and the remain initialized an infinite height, and (2) the water-removing stage, which drains the excess water to ensure that for each cell, there is a path that leads to the boundary [11]. Algorithm 1 is the pseudocode of the P&D algorithm. Line 1 implements the watercovering stage and the intensive computing part of the algorithm is the water-removal stage, an iterative process (lines 2-18).

Algorithm 1 Pseudo code of sequential P&D DEM						
preprocessing algorithm. (zDEM and wDEM are the input						
DEM and the output result respectively, gap is the very small						
slope gradient.)						
viod sDEMProcessing(zDEM, wDEM, gap)						
1: wDEM = Watercovering(zDEM)						
2: do						
3: finished = true						
4: for each c in wDEM do						
5: if $wDEM(c) > zDEM(c)$ then						
6: $\min_neighbor \leftarrow minimum value of neighbors of c$						
7: if $zDEM(c) \ge min_neighbor + gap then$						
8: $wDEM(c) = zDEM(c)$						
9: finished = false						
10: else						
11: if wDEM(c) > min_neighbor + gap then						
12: $wDEM(c) = min_neighbor + gap$						
13: finished = false						
14: end if						

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15: end if
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16: end if
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17: end for
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18: loop while not finished

B. Flow routing

The D8 (deterministic eight-node) algorithm, proposed by O'Callaghan and Mark, directs flow from each grid cell to one of eight neighbors based on slope gradient [4]. The simplest method of calculating flow direction is to determine the slope (Si) to each neighbor and set it to the direction for which Si is the greatest [3]. For each cell, there are two steps to complete the calculation: computing the maximum gradient of eight neighbors and encoding the flow direction to the cell as Figure 1.

a	Z_6	Z_7	Z_8	b	32	64	128
	Z_5	Z _c	Z_1		16		1
	Z_4	Z3	Z ₂		8	4	2

Figure 1. Three-by-three subgrid for the D8 algorithm calculation: (a) Node numbering convention. (b) Flow direction encoding convention.

Once the flow direction is achieved, it could be applied to calculate flow accumulation for each cell. For a cell, the flow accumulation means the flow of all upstream cells which flow into the cell. Assuming each cell has one unit of water, its flow accumulation is the water of all neighbors which drain into it plus their own water (i.e., one multiplies their number):

$$A(c) = n + \sum_{D} A(D) \tag{1}$$

where c denotes any grid cell, n is the number of neighbors flowing into c and D is each neighbor of c whose flow direction is towards c. To calculate the flow accumulation of a cell c, all cells in the region that drain into c must first be calculated [17]. It is means that a cell could achieve its flow accumulation only after all the neighbors which drain into the cell implement their calculation. Assuming m is the number of neighbors that drain into cell c, the flow accumulation of c could be calculated when m equals to 0. Along with m of each cell gradually increased to 0, the flow-accumulation calculation would be done.

The sequential flow-routing algorithm is shown in Algorithm 2. Lines 1-8 serve to achieve the flow-direction grid. In lines 9-20, an iteration process is used for flow accumulation calculation. It is no doubt that the time consuming of flow

accumulation computation is much more than that of flow direction.

Algorithm 2 Pseudo code of serial flow-routing algorithm. (wDEM is the input DEM after preprocessing, FDIR and FACC are the output flow direction and flow accumulation.) viod sFlowRouting (wDEM, FDIR, FACC) 1: for each c in wDEM do 2: if c is not boundary of wDEM then 3: i ← the greatest slope between c and neighbor i 4: FDIR (c) = $2^{(i-1)}$ 5: else 6: FDIR(c) = NODATA7: end if 8: end for 9: do 10: finished = true 11: for each c in FDIR do if flow accumulation of c is not calculated then 12: 13: m ←number of neighbors that drain into cell c 14: if m == 0 then $FACC(c) \leftarrow$ accumulation calculation by Eq. 1 15. 16: finished = false 17: end if 18: end if 19: end for 20: loop while not finished

C. Parallel computation

The first requirement to extract drainage network in parallel is to partition the data among processors. The paper uses a method proposed by Wallis [17], namely, striped partitioning scheme. In this approach, the input dataset is divided horizontally into *size* equal partitions to *size* processors, with any extra portion remaining being attached to the last subdomain. Each process reads in its assigned partition from a grid file, along with an attached buffer which consists of a row of cells as overlapping area from adjacent partitions above and below it. The buffers could be easily updated by obtaining information from adjoining partitions. From these buffers, each processor has access to neighboring partitions with little communication.

According to the above analysis, the parallel algorithms that integrate single program, multiple data (SPMD) are developed. Algorithm 3 and Algorithm 4 are the parallel programs for drainage-network extraction step1 and step 2 respectively, along with Algorithm 5 for step 3. In these parallel algorithms, the function *ShareBuffers()* severs to update the buffers of each partition - the last row of processor P_i is copied to the top buffer of processor P_{i+1} and the top row of processor P_{i+1} is copied to the function *RingTermination()*, the final signal gets 0 if any processor is 0, in addition, the final signal gets 1 while all processors are 1, here, 0 denotes an incompletion status. For Algorithm 3, the function *Waterremoving()* could find its implementation in lines 3-17 of

Algorithm 1. And in line 4 of Algorithm 4, the function *Flowaccumulation()* is completed as lines 10-19 in Algorithm 2.

Alg	orithm 3 Pseudo code of parallel P&D DEM preprocessing
algo	orithm. (zDEM and wDEM are the input DEM and the
out	put result respectively, gap is the very small slope gradient.)
vio	d pDEMProcessing(zDEM, wDEM, gap)
1:	wDEM = Watercovering(zDEM)
2:	do
3:	ShareBuffers(wDEM)
4:	finished = Waterremoving(zDEM, wDEM, gap)
5:	finished = RingTermination(finished)
6٠	loon while not finished
0.	loop white not initiated
d.	arithm A Pseudo code of parallel flow routing algorithm
Alg	orithm 4 Pseudo code of parallel flow-routing algorithm.
Alg (wI FA	orithm 4 Pseudo code of parallel flow-routing algorithm. DEM is the input DEM after preprocessing , FDIR and CC are the output flow direction and flow accumulation)
Alg (wI FA	orithm 4 Pseudo code of parallel flow-routing algorithm. DEM is the input DEM after preprocessing , FDIR and CC are the output flow direction and flow accumulation.) 1 pFlowRouting(wDEM FDIR FACC)
Alg (wI FA vio	orithm 4 Pseudo code of parallel flow-routing algorithm. DEM is the input DEM after preprocessing , FDIR and CC are the output flow direction and flow accumulation.) d pFlowRouting(wDEM, FDIR, FACC) FDIR = Flowdirection(wDEM)
Alg (wI FA vio 1: 2:	orithm 4 Pseudo code of parallel flow-routing algorithm. DEM is the input DEM after preprocessing , FDIR and CC are the output flow direction and flow accumulation.) d pFlowRouting(wDEM, FDIR, FACC) FDIR = Flowdirection(wDEM) do
Alg (wI FA) vio 1: 2: 3:	orithm 4 Pseudo code of parallel flow-routing algorithm. DEM is the input DEM after preprocessing , FDIR and CC are the output flow direction and flow accumulation.) d pFlowRouting(wDEM, FDIR, FACC) FDIR = Flowdirection(wDEM) do ShareBuffers(FACC)
Alg (wI FA) vio 1: 2: 3: 4:	orithm 4 Pseudo code of parallel flow-routing algorithm. DEM is the input DEM after preprocessing , FDIR and CC are the output flow direction and flow accumulation.) d pFlowRouting(wDEM, FDIR, FACC) FDIR = Flowdirection(wDEM) do ShareBuffers(FACC) finished = Flowaccumulation(FDIR, FACC)
Alg (wI FA) vio 1: 2: 3: 4: 5:	orithm 4 Pseudo code of parallel flow-routing algorithm. DEM is the input DEM after preprocessing , FDIR and CC are the output flow direction and flow accumulation.) d pFlowRouting(wDEM, FDIR, FACC) FDIR = Flowdirection(wDEM) do ShareBuffers(FACC) finished = Flowaccumulation(FDIR, FACC) finished = RingTermination(finished)

Algorithm 5 Pseudo code of parallel network-delineating algorithm. (FACC and DNET are the output flow accumulation and drainage network, respectively.)

viod	pNetwork Delineating (FACC, DNET, threshold)
1:	for each c in FACC do
2:	if FACC(c) > threshold then
3:	DNET(c) = 1
4:	else
5:	DNET (c) $= 0$
6:	end if
7:	end for

III. EXPERIMENTS AND RESULTS

The parallel algorithms proposed in this paper has been implemented in C++ using MPI library and tested on a 10-node cluster linked with a Gigabit Ethernet network. These nodes have the same hardware configuration with Intel Core Quad-Core i3-2350 2.3GHz with 4 GB Memory and operating system. For comparison, the sequential algorithms were executed on a single processor within the cluster. Two datasets on a hilly area of Loess Plateau are chosen as study area, the first one consists of a large grid with size 4589 by 6832, and the second dataset consists of a small grid with size 688 by 1025. The time taken to complete the extraction task was measured using the serial algorithms for each dataset. For the parallel algorithms, the time taken to complete the task was measured using a varying number of processors. Figure 2 shows the total execution time of different processors for drainage-network extraction for the small dataset,

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while Figure 3 provides the total execution time for the large dataset.



Figure 2. Time taken to complete drainage-network extraction on the small DEM



Figure 3. Time taken to complete drainage-network extraction on the large DEM

As seen in Figure 2 and Figure 3, the advantage of parallelization for drainage-network extraction is apparent with the crossover at two processors. The computing time within the two processors is obviously shorter than that of the sequential algorithm. Compared with the serial mode, the parallel results of the large and small datasets are almost identical in terms of increased performance speed, even though the data volumes of the large and small datasets are significantly different from each other. With the increase of processor number, the computational time of the parallel implementation decreases greatly with the crossover at eight processors as the increasing of communication time. Fortunately, the proportion of communication is small and the phase scaled well. The parallel algorithm of drainage-network extraction with a large dataset can achieve good performance acceleration. However, the parallel performance does not increase further when the dataset volume achieves a certain size.

IV. CONCLUSIONS

In this paper, we present a set of parallel algorithms for drainage-network extraction from grid terrain datasets. Compared with the existed algorithms, the parallel DEM preprocessing KeyAuthor et al.

algorithm proposed in the paper could not only fill the depressions, but also remove the flat areas. And the parallel flow routing algorithm makes flow accumulation calculation similar to a local algorithm as using the number of upslope-dependence neighbors of a cell. Furthermore, the parallel strategy is relatively simple and can be easily augmented for other similar terrain analysis parameters.

Although this research deals specifically with drainagenetwork extraction, the drainage-basin delineation is a directly extensional concept for further study, and could be easily implemented based on this paper.

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